

L07

Introduction to Min-max Optimization

50.579 Optimization for Machine Learning

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Recap (GANs)

In **Generative Adversarial Networks** (GANs) one would like to solve

$$\min_{\theta} \max_w \mathbb{E}_{x \sim Q} [D_w(x)] - \mathbb{E}_{z \sim F} [D_w(G_{\theta}(z))]$$

- D_w is the discriminator, G_{θ} the generator.
- Q is the data distribution, F say Gaussian (noise)
- D_w might (or not) capture the probability to **classify data point as true!**

- The aforementioned min-max problem is really **hard! Many challenges!**

GANs (Goodfellow et al.)

In their seminal paper, Goodfellow et al. defined the following min-max problem:

$$\min_{\theta} \max_w \mathbb{E}_{x \sim p_{\text{data}}} [\log D_w(x)] + \mathbb{E}_{z \sim p_{\text{noise}}} [\log(1 - D_w(G_{\theta}(z)))]$$

- D_w is the discriminator, G_{θ} the generator.
- p_{data} is the data distribution, p_{noise} say Gaussian (noise).
- D_w captures the probability to **classify data point as true!**
- D is trying to **maximize prob to assign correct label** to both samples from data and from G .

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Lemma (Optimality). For G fixed, the optimal discriminator D has density

$$D_{w^*}(x) = \frac{p_{\text{data}}(x)}{p_{\text{data}}(x) + p_G(x)},$$

where p_G is the *implicit distribution* of the Generator over the data.

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where p_G is the *implicit distribution* of the Generator over the data.

Proof. For fixed G , D is trying to maximize

$$\int_x \log D(x) p_{\text{data}}(x) dx + \int_z \log(1 - D(G(z))) p_{\text{noise}}(z) dz.$$

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The above is nothing but (set $x = G(z)$)

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$$\int_x \log D(x) p_{\text{data}}(x) dx + \int_x \log(1 - D(x)) p_G(x) dx.$$

Finally, observe that function

$$f(y) = a \log y + b \log(1 - y)$$

achieves maximum at $\frac{a}{a+b}$.

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$$\int_x \log D(x)p_{\text{data}}(x)dx + \int_x \log(1 - D(x))p_G(x)dx.$$

Define cost function $C(G)$

$$C(G) := \mathbb{E}_{x \sim p_{\text{data}}} \left[\log \frac{p_{\text{data}}}{p_{\text{data}} + p_G} \right] + \mathbb{E}_{x \sim p_G} \left[\log \frac{p_G}{p_{\text{data}} + p_G} \right].$$

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GANs (Goodfellow et al.)

Theorem (Global solution). *The global minimum of $C(G)$ is achieved if and only if*

$$p_G = p_{data}.$$

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Proof. Observe that for $p_{data} = p_G$ we get that $C(G) = -\log 4$.

Quick recap $\text{KL}(p||q) = \mathbb{E}_{x \sim p} \left[\log \frac{p(x)}{q(x)} \right]$ is **non-negative!**

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Quick recap $\text{KL}(p||q) = \mathbb{E}_{x \sim p} \left[\log \frac{p(x)}{q(x)} \right]$ is **non-negative!**

Finally observe that

$$C(G) = -\log 4 + \text{KL} \left(p_{\text{data}} \parallel \frac{p_{\text{data}} + p_G}{2} \right) + \text{KL} \left(p_G \parallel \frac{p_{\text{data}} + p_G}{2} \right).$$

Min-max Optimization

GANs motivate the study of min-max optimization (in general **harder** than minimization), i.e., for some continuous function f we want to solve

$$\min_{x \in \mathcal{X}} \max_{y \in \mathcal{Y}} f(x, y)$$

Remarks

- Domains are typically **compact**.
- In general the above problem **might not have** a solution.
- There are guarantees when domains are compact and f is **convex-concave**.

Minimax Theorem

Theorem (**Minimax** by John von Neumann). Let $\mathcal{X} \subset \mathbb{R}^n$ and $\mathcal{Y} \subset \mathbb{R}^m$ be *compact convex sets*. If f is a continuous function that is *convex-concave* it holds

$$\min_{x \in \mathcal{X}} \max_{y \in \mathcal{Y}} f(x, y) = \max_{y \in \mathcal{Y}} \min_{x \in \mathcal{X}} f(x, y)$$

Remarks

- Many applications, especially in **Game Theory**.
- If $f = x^T A y$, and the domains are Δ_n, Δ_m it captures classic **zero sum games**
- The above is the **value** of the game.
- Note that It is always true (min-max inequality):

$$\inf_{x \in \mathcal{X}} \sup_{y \in \mathcal{Y}} f(x, y) \geq \sup_{y \in \mathcal{Y}} \inf_{x \in \mathcal{X}} f(x, y)$$

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Define $g(z) \triangleq \inf_{w \in W} f(z, w)$.

$$\forall w, \forall z, g(z) \leq f(z, w)$$

$$\implies \forall w, \sup_z g(z) \leq \sup_z f(z, w)$$

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Proof. Let's use **no-regret learning** for both "players"!

Online Gradient Descent (Recap)

Definition (Online Gradient Descent). Let $f : \mathbb{R}^n \rightarrow \mathbb{R}$ be convex function, differentiable and L -Lipschitz in some compact convex set \mathcal{X} of diameter D . Online GD is defined:

Initialize at some x_0 .

For $t:=1$ to T do

1. Choose x_t and observe $\ell_t(x_t)$.

2. $y_t = x_t - \alpha_t \nabla \ell_t(x_t)$.

3. $x_{t+1} = \Pi_{\mathcal{X}}(y_t)$.

Regret: $\frac{1}{T} \left(\sum_{t=1}^T \ell_t(x_t) - \min_x \sum_{t=1}^T \ell_t(x) \right)$.

Analysis of Online GD for L -Lipschitz (Recap)

Theorem (Online Gradient Descent). Let $f : \mathbb{R}^n \rightarrow \mathbb{R}$ be convex function, differentiable and L -Lipschitz in some compact convex set \mathcal{X} of diameter D . It holds

$$\left(\frac{1}{T} \sum_{t=1}^T \ell_t(x_t) - \min_x \sum_{t=1}^T \ell_t(x) \right) \leq \frac{3}{2} \frac{LD}{\sqrt{T}},$$

with appropriately choosing $\alpha = \frac{D}{L\sqrt{t}}$.

Remarks:

- If we want error ϵ , we need $T = \Theta\left(\frac{L^2 D^2}{\epsilon^2}\right)$ iterations (same as GD for L -Lipschitz).

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Proof. Let's use *no-regret learning* for both "players"!

Let x_1, \dots, x_T and y_1, \dots, y_T be the iterates as advised by some no-regret algorithm and define $\hat{x} = \frac{1}{T} \sum_{i=1}^T x_i$ and $\hat{y} = \frac{1}{T} \sum_{i=1}^T y_i$ and $T = \Theta(\frac{1}{\epsilon^2})$.

Choose any x , then from the *no-regret* property for x we get that

$$\frac{1}{T} \sum_t f(x_t, y_t) \leq \frac{1}{T} \sum_t f(x, y_t) + \epsilon$$

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$$\begin{aligned} \frac{1}{T} \sum_t f(x_t, y_t) &\leq \frac{1}{T} \sum_t f(x, y_t) + \epsilon \\ &\leq f(x, \hat{y}) + \epsilon \text{ by concavity.} \end{aligned}$$

Minimax Theorem

Proof cont.

Choose any y , then from the **no-regret** property for y we get that

$$\begin{aligned}\frac{1}{T} \sum_t f(x_t, y_t) &\geq \frac{1}{T} \sum_t f(x_t, y) - \epsilon \\ &\geq f(\hat{x}, y) - \epsilon \text{ by convexity.}\end{aligned}$$

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We conclude that for all x, y we have

$$f(\hat{x}, y) - 2\epsilon \leq f(x, \hat{y}).$$

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We conclude that for all x, y we have

$$\max_y f(\hat{x}, y) - 2\epsilon \leq \min_x f(x, \hat{y}).$$

Finally we get $\max_y \min_x f(x, y) \geq \min_x f(x, \hat{y})$

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Minimax Theorem

Proof cont.

Choose an

Set $\epsilon \rightarrow 0$ and we are done!

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Last iterate convergence?

Convex-concave settings (with compact domains) are **easy**.
Nevertheless in GANs

- Functions **are not** necessarily convex-concave.
- Time averaging does not help (Jensen's ineq not applicable).
- Motivation to care about **last iterate convergence!**

For the rest of the lecture let's focus on

$$\min_{x \in \mathcal{X}} \max_{y \in \mathcal{Y}} x^T A y.$$

Can we guarantee last iterate convergence using GD or MWUA?

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Can we guarantee last iterate convergence using GD or MWUA?

Not really...

Last iterate convergence

Consider **Gradient Descent/Ascent** that is

$$\begin{aligned}x_{t+1} &= x_t - \eta \nabla_x f(x_t, y_t), \\y_{t+1} &= y_t + \eta \nabla_y f(x_t, y_t).\end{aligned}$$

Consider the simplest case $f(x, y) = xy$. GDA boils down to:

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Claim (Divergence). *It holds that $x_t^2 + y_t^2$ is increasing in t .*

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Claim (Divergence). *It holds that $x_t^2 + y_t^2$ is increasing in t .*

Proof.

$$x_{t+1}^2 + y_{t+1}^2 = (\eta^2 + 1)(x_t^2 + y_t^2).$$

Last iterate convergence

Consider **MWUA** that is

$$x_i^{t+1} = \frac{x_i^t e^{-\eta(Ay^t)_i}}{Z_x},$$

$$y_j^{t+1} = \frac{y_j^t e^{\eta(A^T x^t)_j}}{Z_y}.$$

Theorem (Divergence). *Assume there exists a unique fully mixed Nash (x^*, y^*) equilibrium (full support). It holds that the KL divergence between a player strategies the fully mixed Nash goes to infinity, i.e,*

$$\lim_t \text{KL}(x^* || x^t) = \infty \text{ and } \lim_t \text{KL}(y^* || y^t) = \infty.$$

Conclusion

- Introduction to min-max optimization.
 - GANs.
 - Minimax Theorem
 - Last iterate convergence?
- Next lecture we will talk more about **min-max optimization and optimism.**